

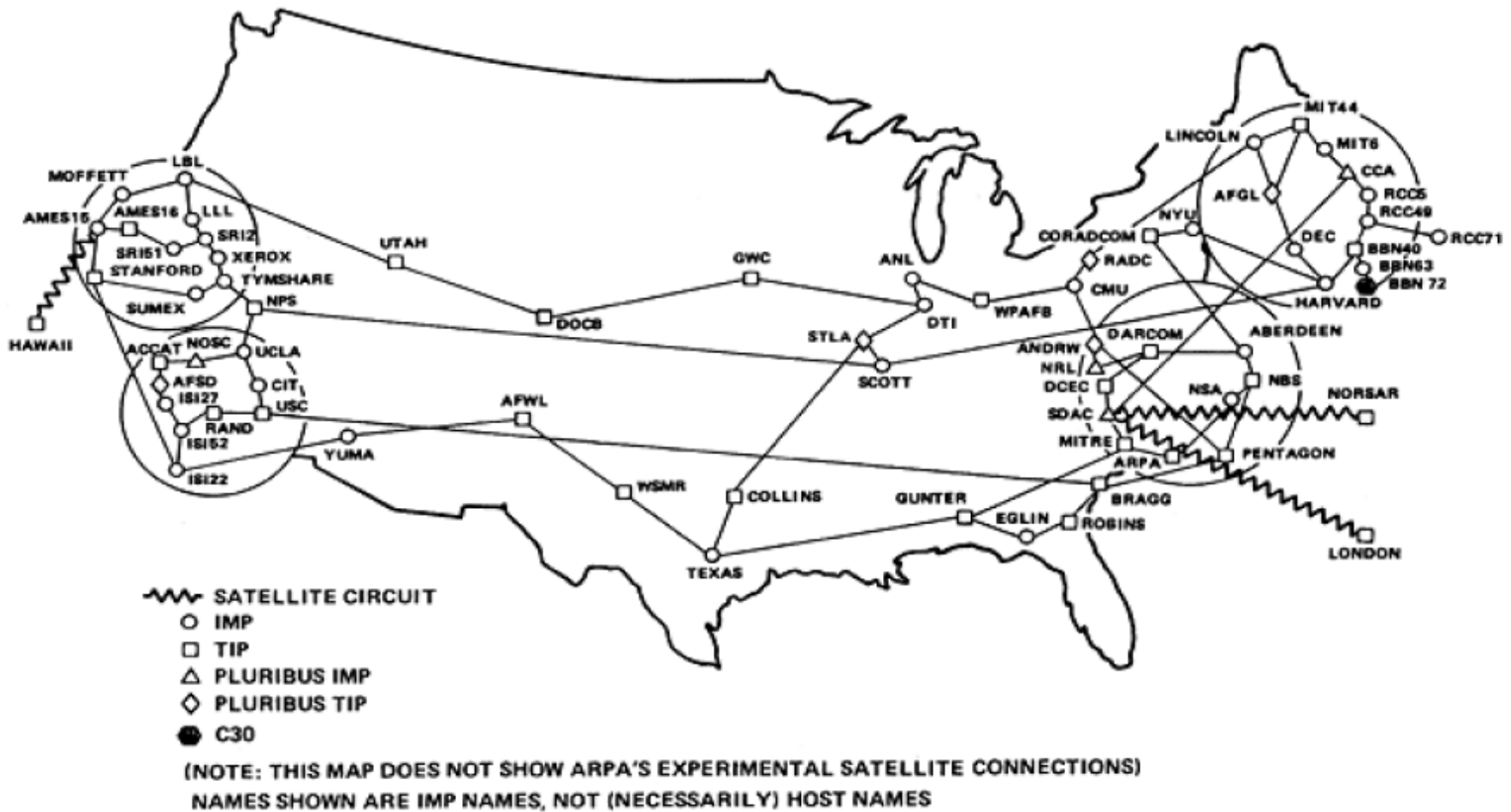
# Anticipation Games

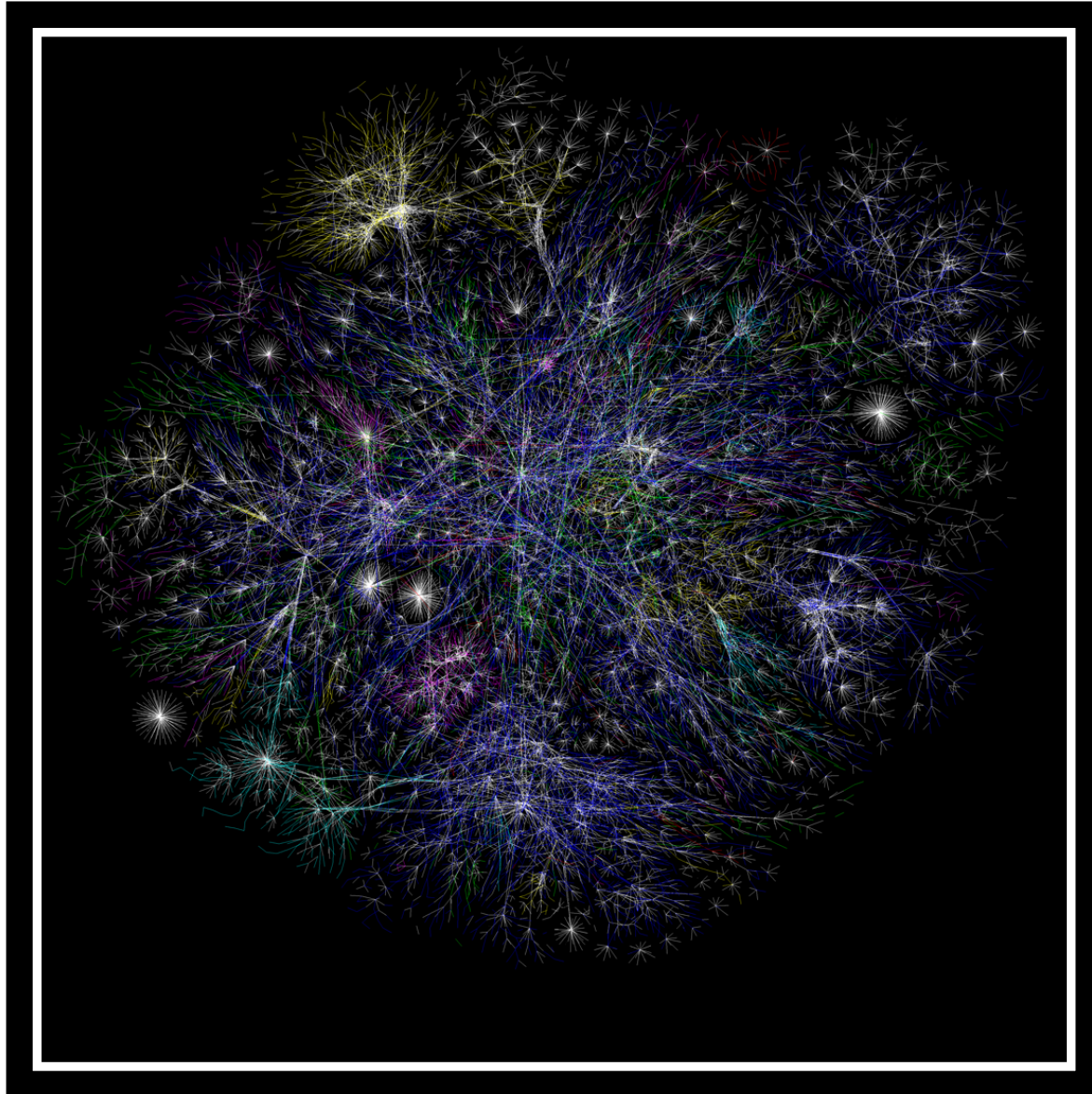
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- ▢ Introduction
  - ▢ Network Evolution
  - ▢ Attack Model Evolution
- ▢ Anticipation game key features
  - ▢ Dependency relations
  - ▢ Player interaction
  - ▢ Time
- ▢ Model Logic
  - ▢ Positional Logic
  - ▢ Temporal Logic
- ▢ Conclusion

ARPANET GEOGRAPHIC MAP, OCTOBER 1980

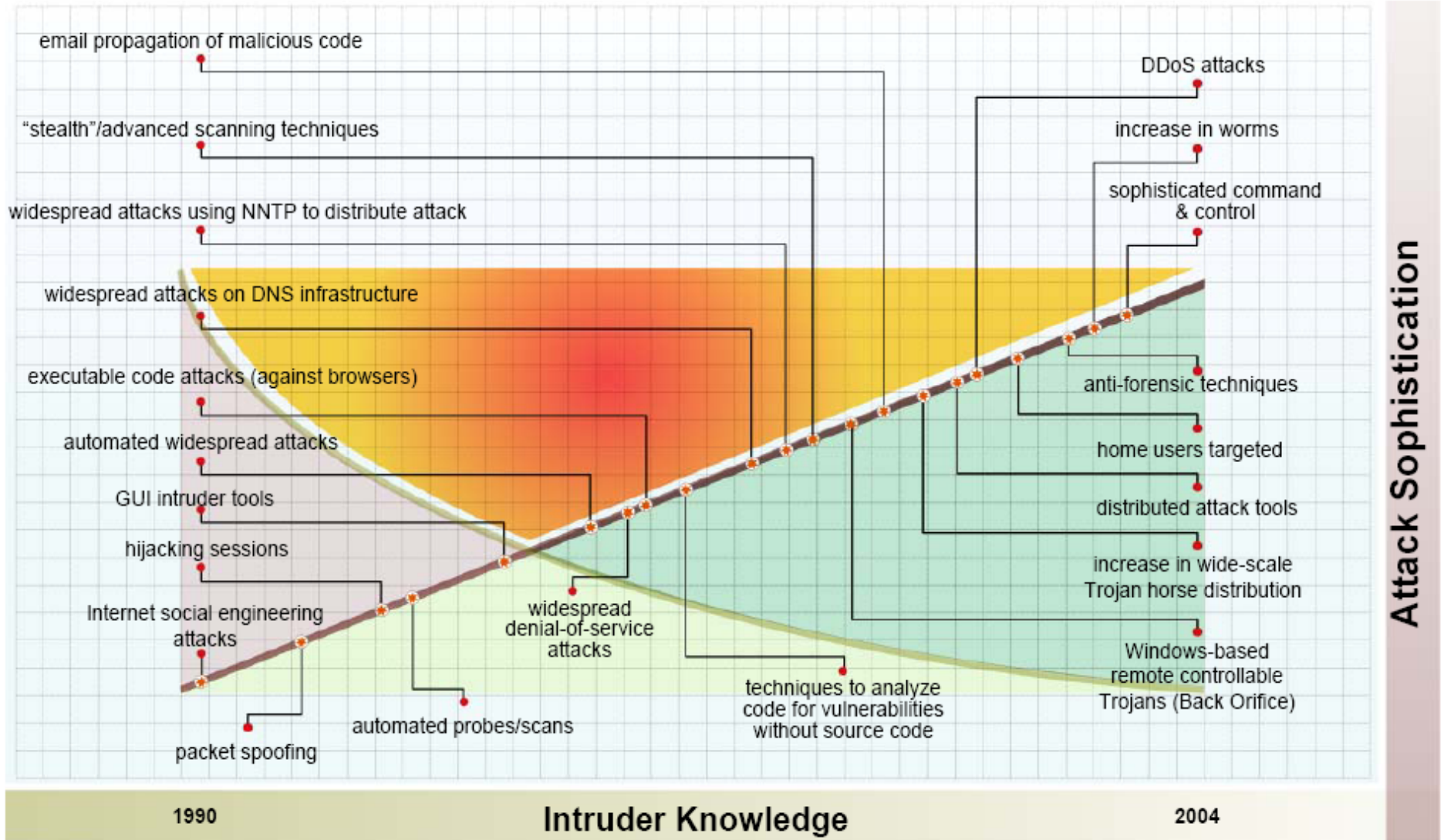




Opte project



# Attack Sophistication vs. Intruder Knowledge



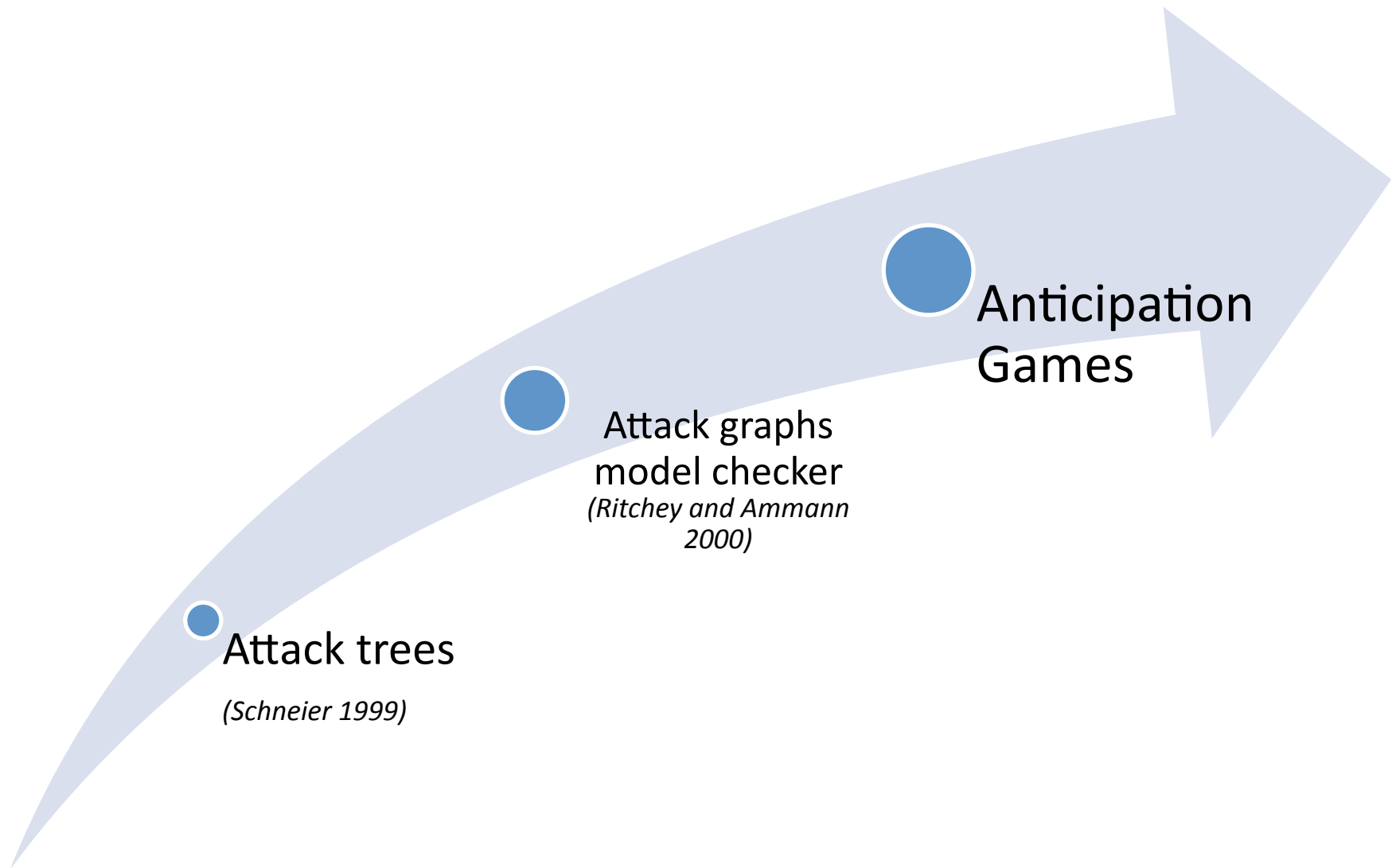
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- Large network may suffer multiple vulnerabilities
- Patches and counter-measures need to be prioritized
- A minor vulnerability can turn into a major hole when used as a step-stone

**Attack graph allows to reason  
about attack sequences**









### Attack graph

- Model checker-based (Ritchey et. al S&P'00, Sheyner et. al S&P'02)
- Graph-based (Ammann et. al CCS'02, Ritchey et. al ACSAC'02, Noel et. al ACSAC'03, Wang et. al ESORICS'05, Wang et. al DBSEC'06)

### Timed Game

- ATL (Alur et al. 97)
- The Element of Surprise in Timed Games (De Alfaro et al. CONCUR 2003)
- TATL (Henzinger et al 2006 Formats)

## Dependency

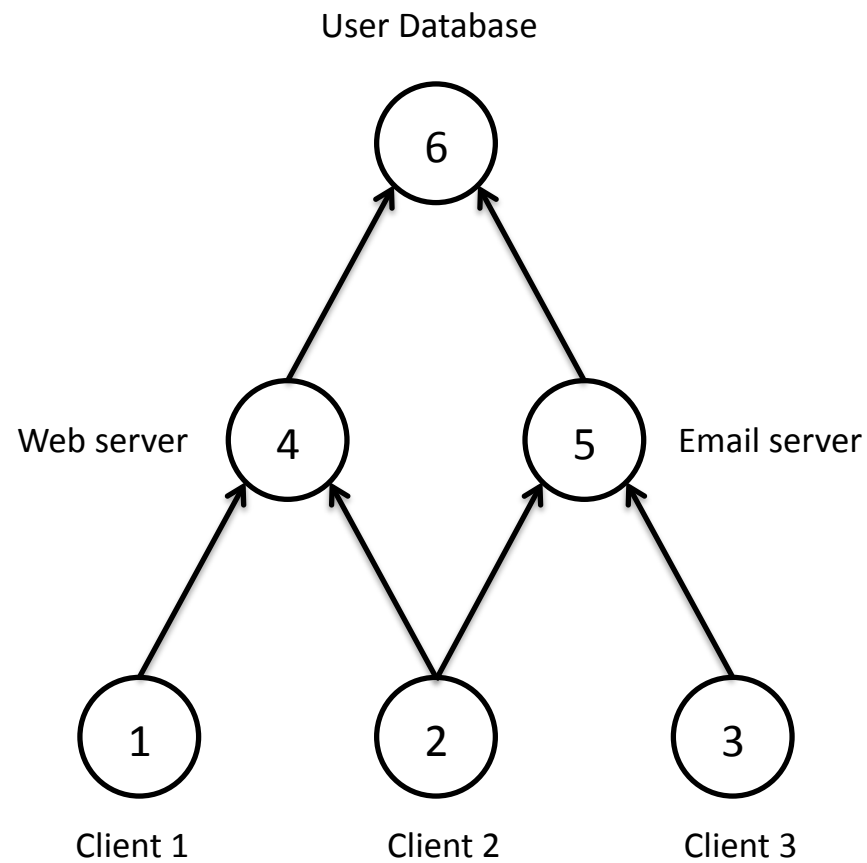
- Collateral effects
- Trust relations

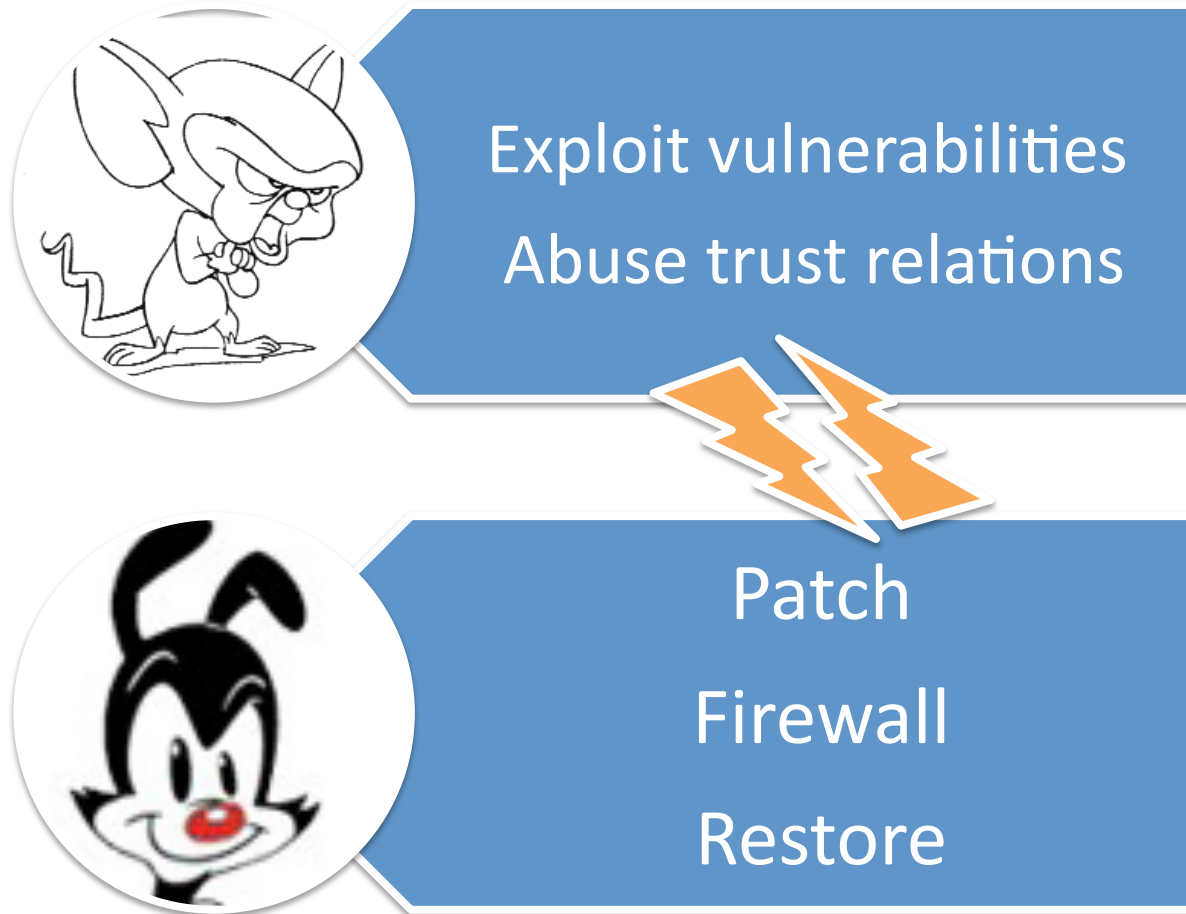
## Interaction

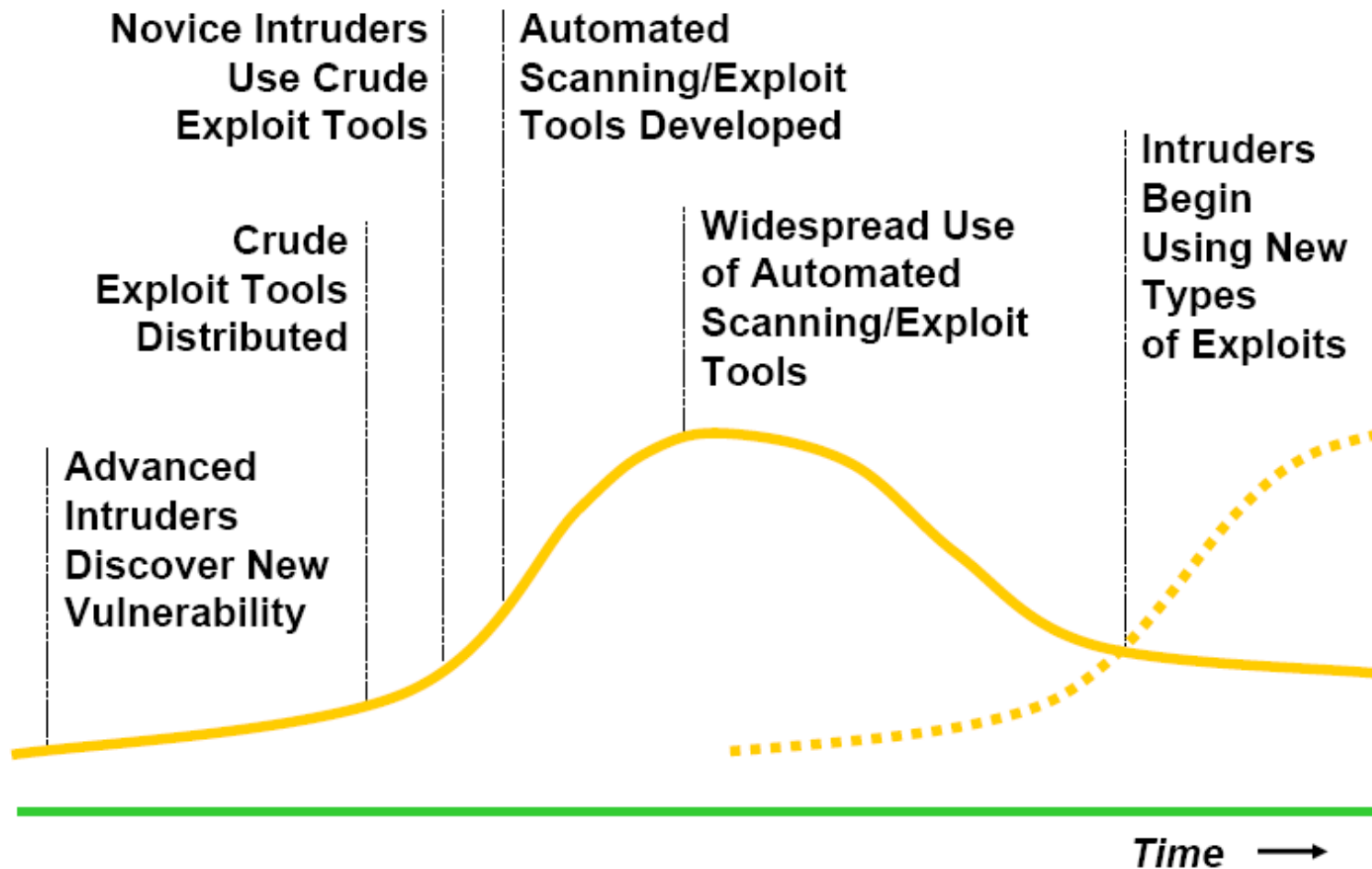
- Administrator
- Intruder

## Time

- Action take time





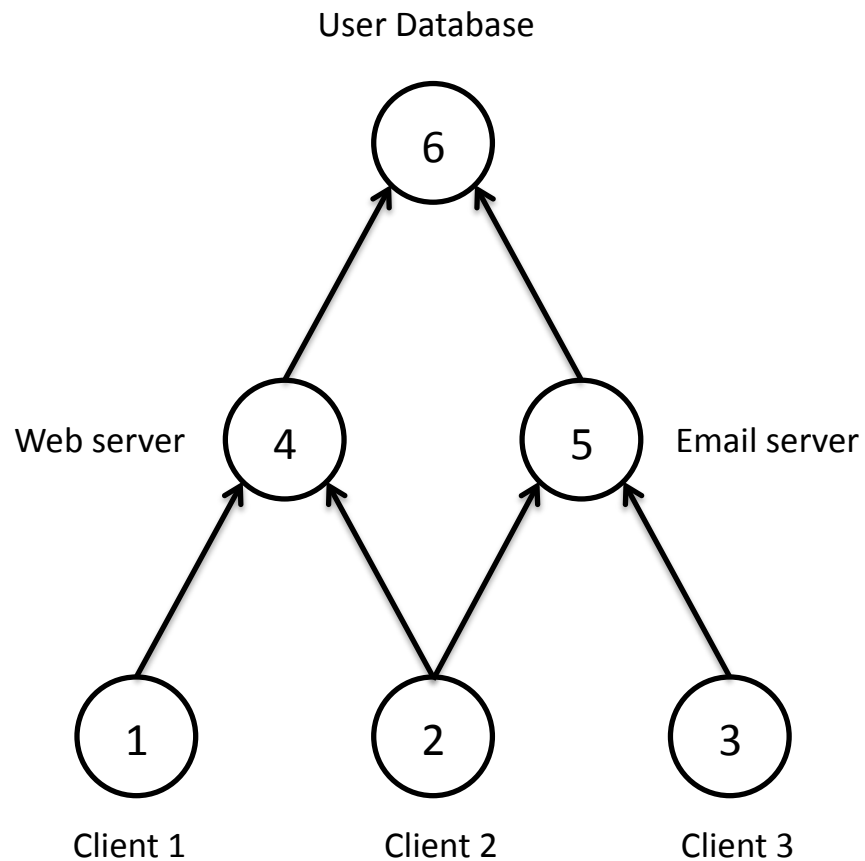


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## Fixed over the time

## Evolve over time



	1	2	3	4	5	6
$\rho(\text{Public})$	⊥	⊥	⊥	T	T	⊥
$\rho(\text{Vuln})$	⊥	⊥	⊥	T	T	⊥
$\rho(\text{Compr})$	⊥	⊥	⊥	⊥	⊥	⊥
$\rho(\text{NeedPub})$	⊥	⊥	⊥	T	T	⊥

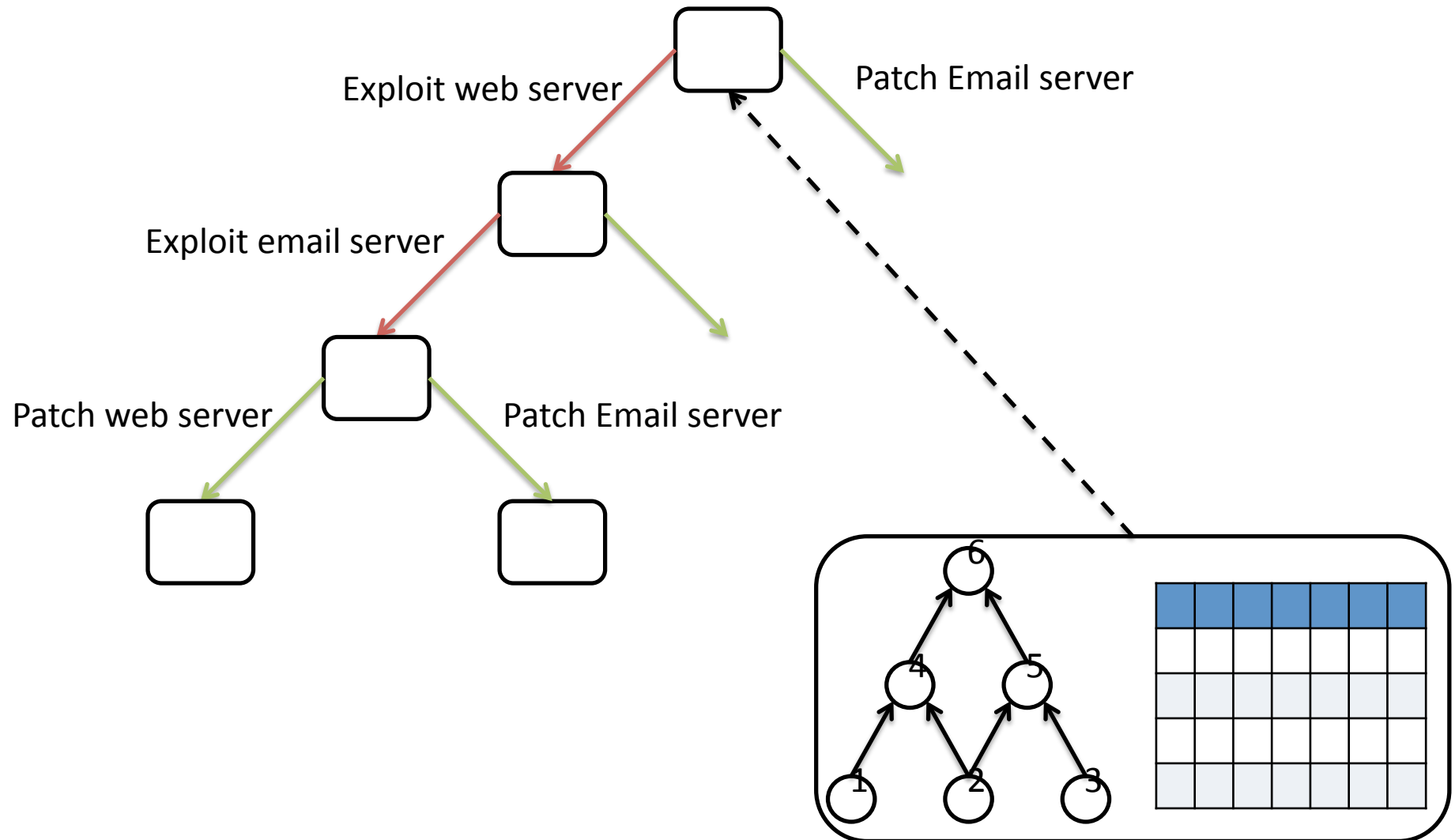
	1	2	3	4	5	6
$\rho(\text{Public})$	$\perp$	$\perp$	$\perp$	T	T	$\perp$
$\rho(\text{Vuln})$	$\perp$	$\perp$	$\perp$	T	T	$\perp$
$\rho(\text{Compr})$	$\perp$	$\perp$	$\perp$	$\perp$	$\perp$	$\perp$
$\rho$ (NeedPub )	$\perp$	$\perp$	$\perp$	T	T	$\perp$

Compr 4

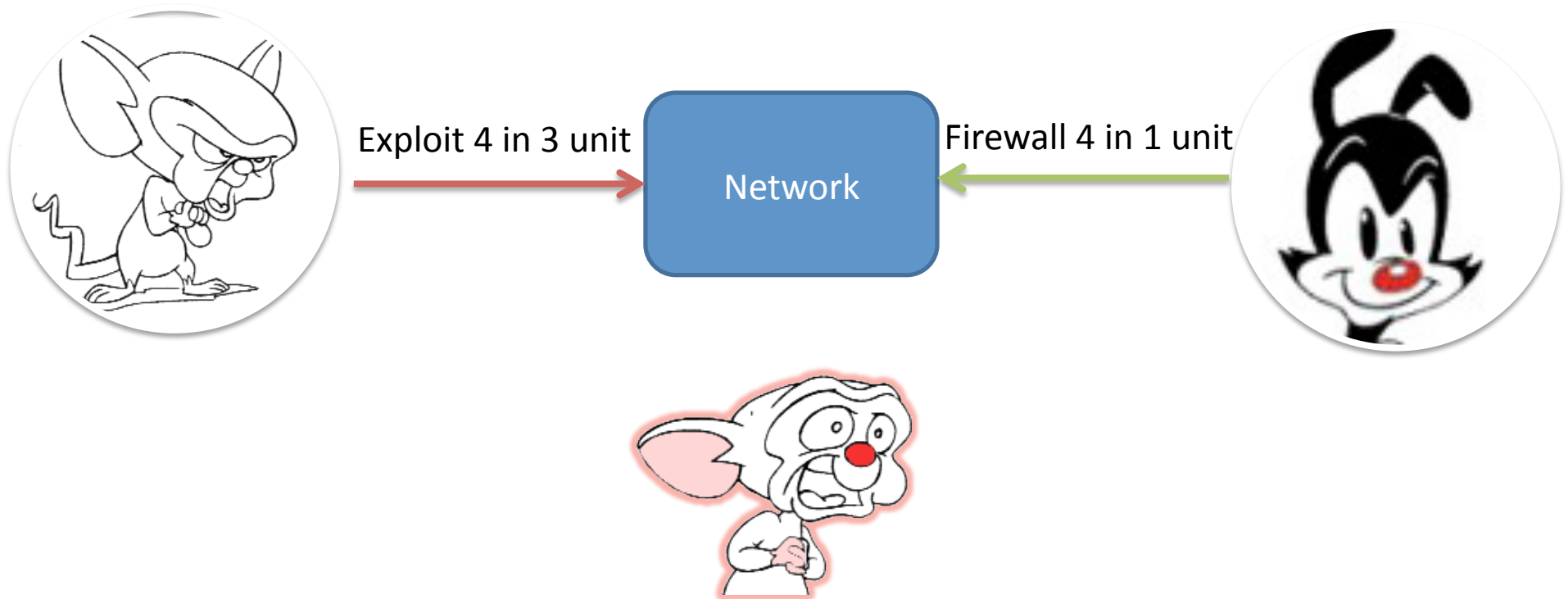


	1	2	3	4	5	6
$\rho(\text{Public})$	$\perp$	$\perp$	$\perp$	T	T	$\perp$
$\rho(\text{Vuln})$	$\perp$	$\perp$	$\perp$	T	T	$\perp$
$\rho(\text{Compr})$	$\perp$	$\perp$	$\perp$	T	$\perp$	$\perp$
$\rho$ (NeedPub )	$\perp$	$\perp$	$\perp$	T	T	$\perp$

# A Incomplete Game Example



- Each action requires a different amount of time
  - Patching a service: Download, extract, apply, restart
  - Exploit a service
  - Firewalling a service
- In anticipation games as in TATL the fastest action win
- Player can be taken by surprise



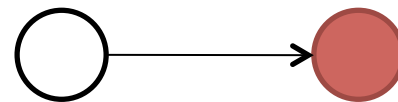


- Anticipation games allows to model
  - Denial of service
  - Buffer overflow execution
  - Permission abuse
  - Cross-scripting
  - Information leak
  - ....

$F$	$::=$	$A$	atomic propositions, in $\mathcal{A}$
		$\top$	true
		$\neg F$	negation
		$F \wedge F$	conjunction
		$\diamond F$	
		$\diamond_{\equiv} F$	

$\vdash \diamond Compr$

A successor node is compromised



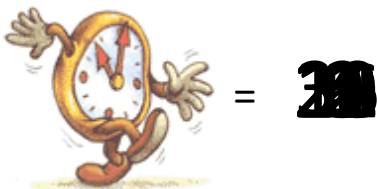
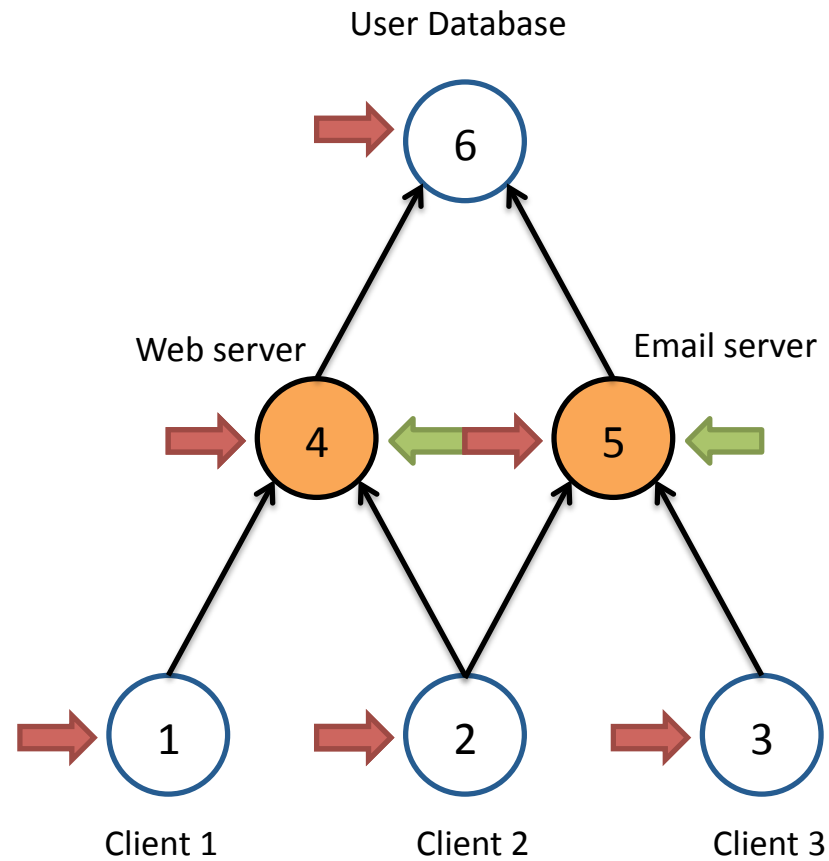
$\vdash \diamond_{\equiv} Public$

At least, one of the node the belongs to the  
equivalence is public



<b>Pre</b> $Vuln \wedge Public \wedge \neg Compr$	$(2, I, \text{Compromise } 0day)$ $\longrightarrow$	$Compr$
<b>Pre</b> $Vuln \wedge Public \wedge \neg Compr$	$(7, I, \text{Compromise } public)$ $\longrightarrow$	$Compr$
<b>Pre</b> $\neg Compr \wedge \diamond Compr$	$(4, I, \text{Compromise } backward)$ $\longrightarrow$	$Compr$
<b>Pre</b> $Compr \wedge \diamond \neg Compr$	$(4, I, \text{Compromise } forward)$ $\longrightarrow$	$\diamond Compr$
<b>Pre</b> $Public \wedge Vuln$	$(1, A, \text{Firewall})$ $\longrightarrow$	$\neg Public$
<b>Pre</b> $Public \wedge \neg Vuln \wedge NeedPub$	$(1, A, \text{UnFirewall})$ $\longrightarrow$	$Public$
<b>Pre</b> $Vuln \wedge \neg Compr$	$(3, A, \text{Patch})$ $\longrightarrow$	$\neg Vuln \wedge \neg Compr$

# A Play example



Player	Action	Rule	Target	Succ
Admin	Execute	Define Wall	5	
Intruder	Execute	Compromise Gateway Backward	5	5



$\varphi$	$::=$	$A$	atomic propositions, in $\mathcal{A}$
		$\neg\varphi$	
		$\varphi \wedge \varphi$	
		$\diamond\varphi$	
		$\diamond\equiv\varphi$	
		$x + d_1 \leq y + d_2$	clock constraints
		$x \cdot \varphi$	freeze
		$\langle\langle\mathcal{P}\rangle\rangle\blacksquare\varphi$	invariant
		$\langle\langle\mathcal{P}\rangle\rangle\varphi_1 \mathcal{U} \varphi_2$	eventually

We abbreviate  $\langle\langle\mathcal{P}\rangle\rangle\text{TRUE} \mathcal{U} \varphi$  as  $\langle\langle\mathcal{P}\rangle\rangle\blacklozenge\varphi$ .

$\vdash \langle\langle A \rangle\rangle \varphi$ 

The player A have a strategy to satisfy  
the property  $\varphi$

 $\vdash \blacksquare \textit{Compr}$ 

In every future the node will be  
compromised

$$\langle\langle A \rangle\rangle \blacksquare \diamond \equiv \neg \text{Compr}$$

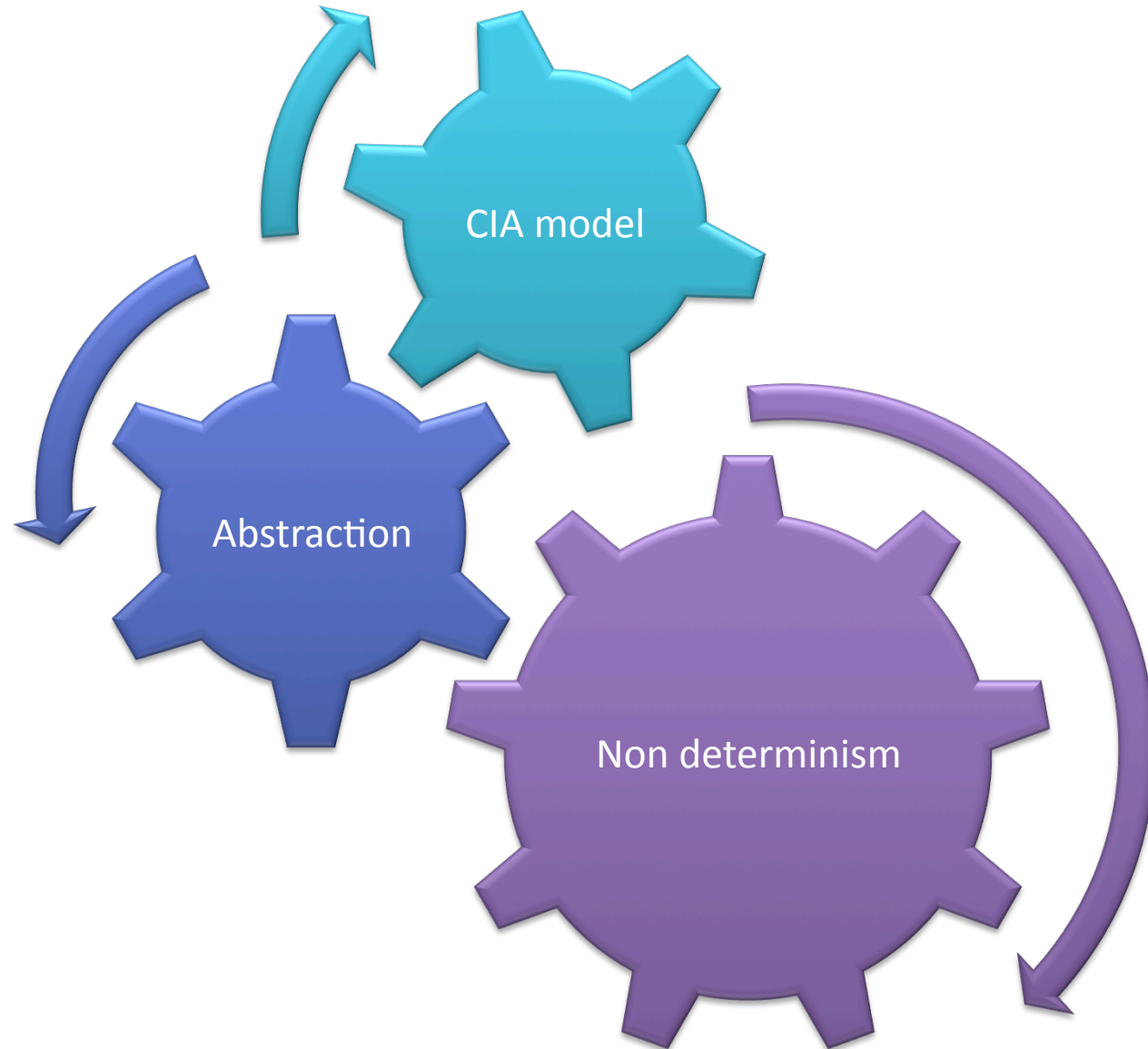
$$\begin{aligned} \langle\langle A \rangle\rangle \blacksquare x \cdot \neg \diamond \equiv \text{Avail} \Rightarrow \\ [\langle\langle A \rangle\rangle \blacklozenge y \cdot y \leq x + d \wedge \langle\langle A \rangle\rangle \blacksquare z \cdot z \leq y + d' \Rightarrow \\ \diamond \equiv \text{Avail}] \end{aligned}$$

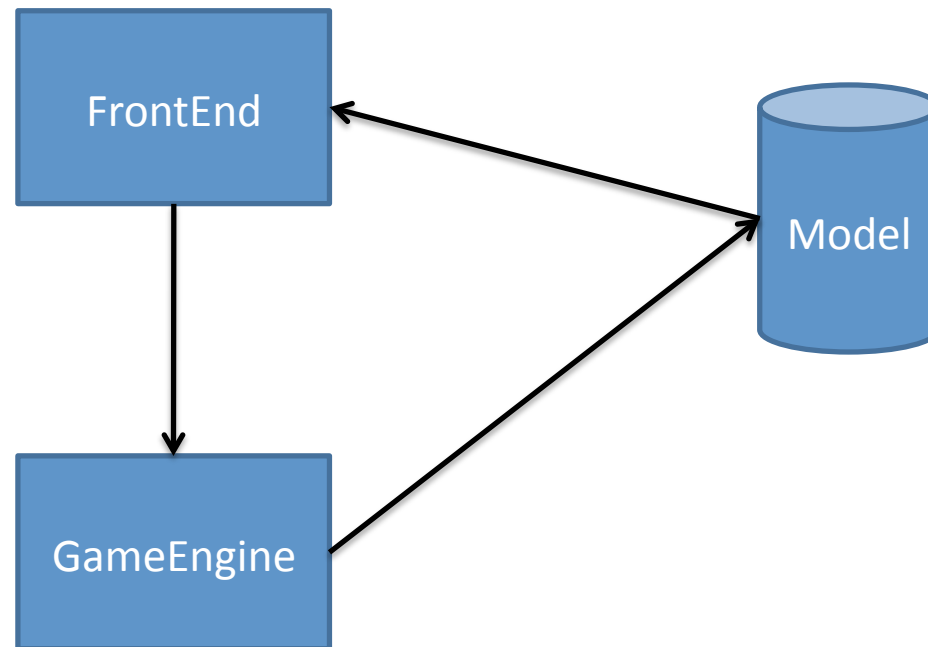
Anticipation game are EXPTIME-complete

One More Thing !

- Model and Strategies are fully implemented in C
- The talk example cannot be analyzed by hand
  - 4011 plays
  - 40825 states









## Analyzer Demo



- Game and Time provide a richer model for intrusion analysis
- Many directions to explore



During this work no network service was injured or tortured.

